

# Disjoint-set Operations

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2019/05/07 @ TR-310-1, NTUST

# Disjoint Sets.

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- Some applications involve grouping  $n$  distinct elements into a collection of disjoint sets
  - Example:
    - $\{1, 2, 3\}$
    - $\{4, 5\}$
    - $\{6\}$
    - $\{1, 3, 5\}$
    - $\{1, 2, 3\}, \{4, 5\}, \{6\}$  are disjoint sets
    - $\{1, 2, 3\}, \{4, 5\}, \{6\}, \{1, 3, 5\}$  are not disjoint sets

# Disjoint Sets..

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- A *disjoint-set data structure* maintains a collection  $S = \{S_1, S_2, \dots, S_k\}$  of disjoint dynamic sets
  - We identify each set by a *representative*
    - Representative should be time-invariant
    - How to choose the representative
- Letting  $x$  denote an object, an element in a set, we wish to support the following operations
  - $\text{MAKE-SET}(x)$  creates a new set whose only member (and thus representative) is  $x$
  - $\text{UNION}(x, y)$  unites the dynamic sets that contain  $x$  and  $y$  into a new set
  - $\text{FIND}(x)$  returns a pointer to the representative of the (unique) set containing  $x$

# Example.

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- One of the many applications of disjoint-set data structures arises in determining the connected components of an undirected graph
  - Given a graph  $G = (V, E)$ , where  $V = \{a, b, c, d, e, f, g, h, i\}$  and  $E = \{(b, d), (e, g), (a, c), (h, i), (a, b), (e, f), (b, c)\}$ 
    - The procedure CONNECTED-COMPONENTS initially places each vertex  $v$  in its own set
    - Then, for each edge  $(u, v)$ , it unites the sets containing  $u$  and  $v$

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CONNECTED-COMPONENTS( $G$ )
1  for each vertex  $v \in G.V$ 
2      MAKE-SET( $v$ )
3  for each edge  $(u, v) \in G.E$ 
4      if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
5          UNION( $u, v$ )
```

# Example..

- Given a graph  $G = (V, E)$ , where  $V = \{a, b, c, d, e, f, g, h, i\}$  and  $E = \{(b, d), (e, g), (a, c), (h, i), (a, b), (e, f), (b, c)\}$

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```

Edge processed	Collection of disjoint sets									
initial sets	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b,d)	{a}	{b,d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e,g)	{a}	{b,d}	{c}		{e,g}	{f}		{h}	{i}	{j}
(a,c)	{a,c}	{b,d}			{e,g}	{f}		{h}	{i}	{j}
(h,i)	{a,c}	{b,d}			{e,g}	{f}		{h,i}		{j}
(a,b)	{a,b,c,d}				{e,g}	{f}		{h,i}		{j}
(e,f)	{a,b,c,d}				{e,f,g}			{h,i}		{j}
(b,c)	{a,b,c,d}				{e,f,g}			{h,i}		{j}

# Example...

- Finally, the procedure SAME-COMPONENT can determine whether two vertices are in the same connected component

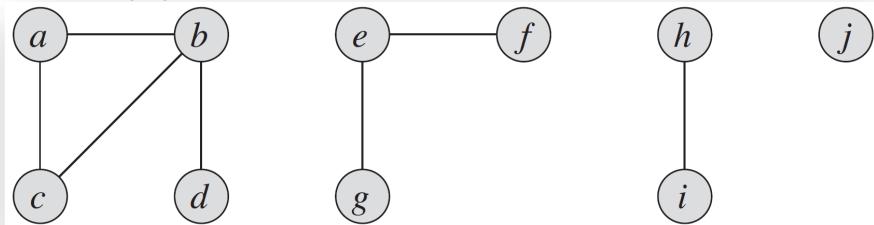
A component, sometimes called a connected component, of an undirected graph is a subgraph in which any two vertices are connected to each other by paths

SAME-COMPONENT( $u, v$ )

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1  if FIND-SET( $u$ ) == FIND-SET( $v$ )
2      return TRUE
3  else return FALSE

```



Edge processed	Collection of disjoint sets									
initial sets	{a}	{b}	{c}	{d}	{e}	{f}	{g}	{h}	{i}	{j}
(b,d)	{a}	{b,d}	{c}		{e}	{f}	{g}	{h}	{i}	{j}
(e,g)	{a}	{b,d}	{c}		{e,g}	{f}		{h}	{i}	{j}
(a,c)	{a,c}	{b,d}			{e,g}	{f}		{h}	{i}	{j}
(h,i)	{a,c}	{b,d}			{e,g}	{f}		{h,i}		{j}
(a,b)	{a,b,c,d}				{e,g}	{f}		{h,i}		{j}
(e,f)	{a,b,c,d}				{e,f,g}			{h,i}		{j}
(b,c)	{a,b,c,d}				{e,f,g}			{h,i}		{j}

# Questions?

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